

A tight lower bound on the minimal dispersion

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Abstract

We give a new lower bound for the minimal dispersion of a point set in the unit cube and its inverse function in the high dimension regime. This is done by considering only a very small class of test boxes, which allows us to reduce bounding the dispersion to a problem in extremal set theory. Specifically, we translate a lower bound on the size of r -cover-free families to a lower bound on the inverse of the minimal dispersion of a point set. The lower bound we obtain matches the recently obtained upper bound on the minimal dispersion up to logarithmic terms.

1 Introduction and the Main Result

For a given subset of the d -dimensional unit cube $X \subset [0, 1]^d$, there are various different ways how to measure whether X is well spread over the unit cube. Based on the previous work of Hlawka [7] and Niederreiter [14], Rote and Tichy [15] introduced in 1996 a concept called the dispersion of X . The *dispersion of X* is the volume of the largest axis-parallel box in $[0, 1]^d$, which does not intersect X , i.e.,

$$\text{disp}(X) = \sup_{B: B \cap X = \emptyset} |B|,$$

where the supremum runs over all the boxes $B = \prod_{i=1}^d (a_i, b_i)$, where $0 \leq a_i < b_i \leq 1$ for all $i \in \{1, \dots, d\}$, and $|B|$ stands for the volume of B .

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Intuitively, the smaller the dispersion of a point set is, the better spread over the unit cube the points have. We are therefore interested in point sets that have simultaneously small cardinality and small dispersion.

Given $n \in \mathbb{N}$, we denote the minimal dispersion of a point set that has cardinality n by

$$\text{disp}^*(n, d) = \inf_{\substack{X \subset [0,1]^d \\ \#X=n}} \text{disp}(X).$$

For a fixed $d \in \mathbb{N}$ and $\varepsilon \in (0, 1)$, we define the inverse function of the minimal dispersion in $[0, 1]^d$ by

$$\begin{aligned} N(\varepsilon, d) &= \min\{n \in \mathbb{N} : \text{disp}^*(n, d) \leq \varepsilon\} \\ &= \min\{n \in \mathbb{N} : \exists X \subset [0, 1]^d \text{ with } \#X = n \text{ and } \text{disp}(X) \leq \varepsilon\}. \end{aligned}$$

The study of the minimal dispersion $\text{disp}^*(n, d)$ as well as its inverse function $N(\varepsilon, d)$ recently attracted a remarkable attention in the mathematical community. The pigeonhole principle implies the trivial lower bound $\text{disp}^*(n, d) \geq \frac{1}{n+1}$, which translates into

$$N(\varepsilon, d) \geq \frac{1}{\varepsilon} - 1.$$

This bound was complemented in [15] and then further improved by Bukh and Chao [4] to

$$N(\varepsilon, d) \leq C \cdot \frac{d^2 \log d}{\varepsilon}. \quad (1)$$

Let us note that further upper and lower bounds for different regimes of ε and d were obtained also in [5, 11, 12, 13, 16], and the dispersion of certain specific sets was studied in [3, 6, 8, 9, 10, 19, 20].

The first non-trivial lower bound on $N(\varepsilon, d)$ that exhibits a growth if d increases was achieved by Aistleitner, Hinrichs and Rudolf [1]. Specifically, they showed that

$$N(\varepsilon, d) \geq \frac{\log_2 d}{8\varepsilon} \quad (2)$$

for every $d \geq 2$ and every $0 < \varepsilon < 1/4$. The results mentioned so far suggest that $N(\varepsilon, d)$ depends linearly on $1/\varepsilon$ and it only remains to determine the d -dependence.

In another line of research, Sosnovec [18] showed that $N(\varepsilon, d) \leq C_\varepsilon \log d$. The dependence of C_ε was not optimized in [18], and a brief inspection of the proof reveals that it is actually super-exponential in $1/\varepsilon$. Nevertheless, it matches (2) when it comes to the dependence on d . Let us mention that

the point set with small dispersion was generated in [18] by sampling the coordinates of the points independently at random from a uniform grid in $[0, 1]$. An improved analysis of the same method was given in [21], which together with a further improvement due to Litvak [11] yields

$$N(\varepsilon, d) \leq C \cdot \frac{\log d \cdot \log(\frac{1}{\varepsilon})}{\varepsilon^2}. \quad (3)$$

The comparison between (1) and (3) is not easy. Depending on the relation between d and $1/\varepsilon$, one or the other might perform better. Although the analysis in [11, 21] regarding the dependence on $1/\varepsilon$ was much finer than the one in [18], it was quite natural to assume that the second power of $1/\varepsilon$ in (3) is an artefact of the proof method and that the bound could be potentially further improved even when ε is moderately large.

The main result of this work is to show that, surprisingly, this is not the case and the second power of $1/\varepsilon$ in (3) is optimal when ε is sufficiently large.

Theorem 1. *There is an absolute constant $c > 0$, such that the following statement is true. For any integer $d \geq 2$ and any real ε satisfying $\frac{1}{4} \geq \varepsilon \geq \frac{1}{4\sqrt{d}}$, it holds that*

$$N(\varepsilon, d) > \frac{c \log d}{\varepsilon^2 \cdot \log \frac{1}{\varepsilon}}. \quad (4)$$

The proof of this result is motivated by the insight gained by a detailed inspection of the proof given in [21]. For a fixed box B , the sampling from a finite grid, which is bounded away of zero and one, naturally splits the individual coordinates into two groups — those, where the corresponding interval (a_i, b_i) is large enough and covers the whole grid, and the coordinates, where this is not the case. It turned out in [21] that the boxes with a large number of the coordinates inside the second group are the most difficult to hit by a randomly generated point. Therefore, we choose a very small class of test boxes when compared to the number of all the axis-parallel boxes, and consider any point set that hits all these test boxes. This naturally reduces the question under study to a well-known combinatorial problem called *r-cover-free systems*, see Section 2 for the definition. The key tool in our proof is a lower bound of Alon and Asodi [2] on the minimum size of *r-cover-free system*.

Reformulated in the language of minimal dispersion, Theorem 1 states the following.

Corollary 2. *There are two absolute constants $c_1, c_2 > 0$ such that for every positive integers d and n with $d \geq 2$ and $2 \log d \leq n \leq c_1 d$ it holds that*

$$\text{disp}^*(n, d) \geq c_2 \left(\frac{\log d}{n} \right)^{1/2} \cdot \left(\log \left(\frac{n}{\log d} \right) \right)^{-1/2}.$$

Before we present our proof of Theorem 1, let us add a few remarks.

- Remark.**
1. *It follows from the proof of Theorem 1 that one can actually choose $c = 1/1920$, but we did not attempt to optimize this constant.*
 2. *The essentially tight factor $1/\varepsilon^2$ in the lower bound for $N(\varepsilon, d)$ in (4) is rather surprising. It seems that a logarithmic dependence on d and a linear dependence on $1/\varepsilon$ exclude each other. What might be even more surprising, this lower bound is obtained by considering only a very small class of axis-parallel test boxes.*
 3. *We leave it as an open problem whether the proof method can be further improved, possibly extending (4) to smaller values of ε .*

2 Proof of Theorem 1

We start by fixing some notation that will be used in the proof. For a finite set X , we denote its size by $\#X$. For a positive integer d , we denote by $[d]$ the set $\{1, 2, \dots, d\}$. Additionally, for a non-negative integer k , we write $\binom{[d]}{k}$ to denote the collection of all the k -element subsets of $[d]$. Given a point $x \in [0, 1]^d$ and an integer $i \in [d]$, we denote by $(x)_i$ the i -th coordinate of x .

By focusing on a very specific set of axis-parallel boxes of volume at least ε , we are able to translate bounding $N(\varepsilon, d)$ from below to a question in extremal set theory. We say that a family \mathcal{F} of subsets of a ground set X is r -cover-free if no $F \in \mathcal{F}$ is contained in the union of any r sets from $\mathcal{F} \setminus \{F\}$. Based on a previous work of Ruszinkó [17], Alon and Asodi [2] proved the following lower-bound on the size of the ground set of an r -cover-free family.

Theorem 3 ([2, Lemma 2.8]). *Let \mathcal{F} be a family of d subsets of a ground set X . If \mathcal{F} is r -cover-free, where $2 \leq r \leq 2\sqrt{d}$, then*

$$\#X > \frac{1}{10} \cdot \frac{r^2 \log \left(d - \frac{r}{2} \right)}{\log r}.$$

We now proceed with the proof of Theorem 1. Fix $d \geq 2$ and ε such that $\frac{1}{4} \geq \varepsilon \geq \frac{1}{4\sqrt{d}}$, and let k be the unique integer satisfying $2^{-(k+1)} < \varepsilon \leq 2^{-k}$. Note that $k \geq 2$.

Our plan is to define a collection \mathcal{B} of $\binom{d}{2^{k-2}}$ ($d - 2^{k-2}$) boxes such that each box has one coordinate where all its points are close to zero, while in some other 2^{k-2} coordinates all the points of the box are bounded away from zero. Next, we will sort out the points of X with at least one coordinate close to zero into (not necessarily disjoint) sets F_1, F_2, \dots, F_d . All of this is done in such a way that the sets F_1, F_2, \dots, F_d are 2^{k-2} -cover-free, thus Theorem 3 yields the desired lower bound on X .

We now give all the details. For a given set $A \subseteq [d]$ and $j \in [d] \setminus A$, let $B^{A,j} \subseteq [0, 1]^d$ be the axes-parallel box defined as

$$B^{A,j} := I_1 \times I_2 \times \dots \times I_d, \quad \text{where } \begin{cases} I_j = (0, 2^{1-k}), \\ I_i = (2^{1-k}, 1) \quad \text{for } i \in A, \\ I_i = (0, 1) \quad \text{for } i \in [d] \setminus (A \cup \{j\}). \end{cases}$$

Let $\mathcal{B} := \{B^{A,j} : A \in \binom{[d]}{2^{k-2}} \text{ and } j \in [d] \setminus A\}$. The following claim yields that any set of points from $[0, 1]^d$ with dispersion at most ε must have at least one point inside every $B \in \mathcal{B}$.

Claim 1. $|B^{A,j}| \geq \varepsilon$ for any $A \subseteq [d]$ of size at most 2^{k-2} and any $j \in [d] \setminus A$.

Proof. Since 2^{-2x} is convex, it holds that $(1-x) \geq 2^{-2x}$ for every $x \in [0, 1/2]$. Therefore,

$$|B^{A,j}| = \left(1 - 2^{1-k}\right)^{\#A} \cdot 2^{1-k} \geq \left(2^{-2^{2-k}}\right)^{\#A} \cdot 2^{1-k} \geq 2^{-k} \geq \varepsilon,$$

which finishes the proof of the claim. \square

Fix a finite set of points $X \subseteq [0, 1]^d$ that intersects every $B \in \mathcal{B}$. For every $j \in [d]$, we define an auxiliary set F_j of the points with a small j -th coordinate. Specifically, let $F_j := \{x \in X : (x)_j < 2^{1-k}\}$. In order to lower bound $\#X$ using Theorem 3, we establish the following claim.

Claim 2. The family $\{F_1, F_2, \dots, F_d\}$ is 2^{k-2} -cover-free.

Proof. Suppose there is a set $A \in \binom{[d]}{2^{k-2}}$ and an integer $j \in [d] \setminus A$ such that

$$F_j \subseteq \bigcup_{i \in A} F_i.$$

However, there must exist a point $x \in X \cap B^{A,j}$. Since $(x)_j < 2^{1-k}$ and $(x)_i > 2^{1-k}$ for all $i \in A$, we conclude that x is an element of $F_j \setminus \bigcup_{i \in A} F_i$, which is a contradiction. \square

If $k = 2$, then $\varepsilon \in (1/8, 1/4]$ and $\#X \geq \log_2(d) > \frac{3}{64} \cdot \frac{\log d}{\varepsilon^2 \cdot \log(1/\varepsilon)}$ by the 1-cover-freeness. For $k \geq 3$, we have that $2^{k-2} \leq \frac{1}{4\varepsilon} \leq \sqrt{d}$, thus Theorem 3 applied to the family $\{F_1, \dots, F_d\}$ and $r = 2^{k-2}$ readily yields that

$$\#X > \frac{1}{10} \cdot \frac{2^{2k-4} \cdot \log(d - 2^{k-3})}{\log(2^{k-2})} \geq \frac{1}{10} \cdot \frac{2^{2k-4} \cdot \log(d - \sqrt{d}/2)}{\log(2^{k-2})}.$$

Since $\frac{1}{8\varepsilon} < 2^{k-2} < \frac{1}{\varepsilon}$ and $\log(d - \sqrt{d}/2) \geq \frac{\log d}{3}$ for all $d \geq 2$, we conclude that

$$\#X > \frac{1}{1920} \cdot \frac{\log d}{\varepsilon^2 \cdot \log \frac{1}{\varepsilon}}.$$

This finishes the proof of Theorem 1.

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